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# The $r$-acyclic chromatic number of planar graphs 

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#### Abstract

A vertex coloring of a graph $G$ is $r$-acyclic if it is a proper vertex coloring such that every cycle $C$ receives at least $\min \{|C|, r\}$ colors. The $r$-acyclic chromatic number $a_{r}(G)$ of $G$ is the least number of colors in an $r$-acyclic coloring of $G$. Let $G$ be a planar graph. By Four Color Theorem, we know that $a_{1}(G)=a_{2}(G)=\chi(G) \leq 4$, where $\chi(G)$ is the chromatic number of $G$. Borodin proved that $a_{3}(G) \leq 5$. However when $r \geq 4$, the $r$-acyclic chromatic number of a class of graphs may not be bounded by a constant number. For example, $a_{4}\left(K_{2, n}\right)=n+2=\Delta\left(K_{2, n}\right)+2$ for $n \geq 2$, where $K_{2, n}$ is a complete bipartite (planar) graph. In this paper, we give a sufficient


[^0]condition for $a_{r}(G) \leq r$ when $G$ is a planar graph. In precise, we show that if $r \geq 4$ and $G$ is a planar graph with $g(G) \geq \frac{10 r-4}{3}$, then $a_{r}(G) \leq r$. In addition, we discuss the 4 -acyclic colorings of some special planar graphs.

Keywords Acyclic coloring • Planar graph • Girth

## 1 Introduction and notation

We use Bondy and Murty (1976) for terminology and notations not defined here and consider undirected graphs only. Let $G=(V, E)$ be a graph. A vertex coloring of a graph $G$ is $r$-acyclic if it is a proper vertex coloring such that every cycle $C$ receives at least $\min \{|C|, r\}$ colors. The $r$-acyclic chromatic number of $G, a_{r}(G)$, is the least number of colors in an $r$-acyclic coloring of $G$.

For $r \leq 2$, the $r$-acyclic coloring is actually the proper vertex coloring, so for any graph $G$ with maximum degree $\Delta$, its $r$-acyclic chromatic number is at most $\Delta+1$. The 3-acyclic coloring, which is also known as acyclic coloring in the literature, has been studied extensively. It was proved by Skulrattanakulchai (2004) that $a_{3}(G) \leq 4$ for any graph of maximum degree 3 . Burnstein (1979) showed that $a_{3}(G) \leq 5$ for any graph of maximum degree 4. Kostochka and Stocker (2011) proved that $a_{3}(G) \leq 7$ for any graph of maximum degree 5. Hocquard (2011) confirmed that $a_{3}(G) \leq 11$ for any graph of maximum degree 6 . Dieng et al. (2010) showed for each graph $G$ with maximum degree $\Delta \geq 7$ that $a_{3}(G) \leq f(\Delta)$, where

$$
f(\Delta)= \begin{cases}17 & \text { if } \Delta=7  \tag{1.1}\\ \frac{\Delta^{2}-5 \Delta}{2}+2 \times\left[\frac{\Delta-1}{2}\right]+3 & \text { if } \Delta \geq 8\end{cases}
$$

Yadav et al (2009) showed for any graph $G$ with maximum degree $\Delta$ that $a_{3}(G) \leq$ $\frac{3 \Delta^{2}+4 \Delta+8}{8}$. Alon et al. (1991) gave upper and lower bounds for $a_{3}(G)$ by using the probabilistic method; they proved that for some constants $c_{1}, c_{2}>0$,

$$
\frac{c_{1} \Delta^{\frac{4}{3}}}{(\log \Delta)^{\frac{1}{3}}} \leq a_{3}(G) \leq c_{2} \Delta^{\frac{4}{3}}
$$

For $r \geq 4$, it was shown in Greenhill and Pikhurko (2005) that there exist positive constants $c, c^{\prime}$ such that $c \Delta^{\left\lfloor\frac{r}{2}\right\rfloor} \leq a_{r}(G) \leq c^{\prime} \Delta^{\left\lfloor\frac{r}{2}\right\rfloor}$. Cai et al. (2013) proved that for a graph $G$ with maximum degree $\Delta$ and girth $g \geq 2(r-1) \Delta$, $a_{r}(G) \leq 6(r-1) \Delta$, where $r \geq 4$ is a positive integer. For more references, we refer to (Albertson and Berman 1976; Fertin and Raspaud 2005, 2008; Zhang et al. 2012).

Now we focus on planar graphs. First of all, the Four Color Theorem implies $a_{1}(G) \leq 4$ and $a_{2}(G) \leq 4$. (Grünbaum 1973) conjectured that 5 colors are sufficient to acyclically color any planar graph; this conjecture was confirmed by Borodin (1979). However, for $r \geq 4$ and a class of graphs $\mathcal{G}, a_{r}(\mathcal{G}):=\max \left\{a_{r}(G) \mid G \in \mathcal{G}\right\}$ may not be bounded by a constant number, the class of planar complete bipartite graphs is such an example, since $a_{r}\left(K_{2, n}\right)=n+2$ is dependent of the maximum degree, where $r \geq 4$ and $n \geq 2$.

In this paper, we consider $r$-acyclic colorings of planar graphs. First, we discuss the 4 -acyclic colorings of some special planar graphs. Whereafter, we give a sufficient condition for $a_{r}(G) \leq r$ for planar graphs.

## 24-acyclic colorings of outerplanar graphs

A graph is outerplanar if it can be drawn in the plane so that all vertices are lying on the outside face. It is known that any outerplanar graphs contains no $K_{2,3}$-minors and $K_{4}$-minors. In this and the next section, multiple edges are allowed. The following structural lemma for outerplanar graphs proved by Borodin and Woodall (1995) is a useful start.

Lemma 2.1 Borodin and Woodall (1995) Every outerplanar graph with minimum degree at least two contains one of the following configurations:
(a) two adjacent 2-vertices $u$ and $v$;
(b) a 3-cycle uvw with $d(u)=2$ and $d(v)=3$;
(c) two intersecting 3-cycles uvw and $x v y$ with $d(u)=d(x)=2$ and $d(v)=4$.

Theorem 2.1 For each outerplanar graph $G, a_{4}(G) \leq 4$ and the bound 4 is sharp.
Proof Suppose that $G$ is a minimal counterexample with the smallest number of $|V(G)|+|E(G)|$. Clearly, $\delta(G) \geq 2$. By Lemma 2.1, we need consider two cases.

First, suppose that $G$ contains two adjacent 2 -vertices $u$ and $v$. Denote the other neighbor of $u$ and $v$ by $w$ and $z$, respectively. If $z w \in E(G)$, then let $H:=G-\{u, v\}$; otherwise, let $H:=G-\{u, v\}+z w$. In each case one can check that $H$ is still outerplanar and thus by the minimality of $G, H$ admits a 4-acyclic coloring $\varphi$ with $\varphi(z) \neq \varphi(w)$. Extend $\varphi$ to a coloring of $G$ by coloring $u$ and $v$ with two distinct colors that are different from $\varphi(z)$ and $\varphi(w)$. Since every cycle of length at least 4 in $G$ passing through $u$ and $v$ contains the path $w u v z$, the resulting coloring of $G$ is 4-acyclic as required.

Second, suppose that $G$ contains a 2-vertex $u$ that is incident with a triangle $u v w$. Let $H:=G-u$. By the minimality of $G, H$ admits a 4-acyclic coloring $\varphi$. Since $G$ is outerplanar, $|N(v) \cap N(w)| \leq 2$, because otherwise one can find a $K_{2,3}$-minor in $G$. This implies that $u$ is incident with at most one 4-cycle in $G$. If $N(v) \cap N(w)=\{u, s\}$, then extend $\varphi$ to a coloring of $G$ by coloring $u$ with a color different from $\varphi(v), \varphi(w)$ and $\varphi(s)$. If $C$ is a cycle of length at least 5 in $G$ passing through $u$, then $C-u+v w$ is a cycle of length at least 4 in $H$. Since $H$ has already been 4 -acyclically colored, the cycle $C-u+v w$ in $H$ is incident with at least 4 colors under $\varphi$. Thus, the cycle $C$ in $G$ is also incident with at least 4 colors after the extension of $\varphi$. By the choice of the coloring on $u$, the vertices in the unique 4 -cycle that passes through $u$ are colored all distinctly. Therefore, the extended coloring of $G$ is 4 -acyclic, a contradiction. If $N(v) \cap N(w)=\{u\}$, then extend $\varphi$ to a coloring of $G$ by coloring $u$ with a color different from $\varphi(v)$ and $\varphi(w)$. If $C$ is a cycle of length at least 4 in $G$ passing through $u$, then $C-u+v w$ is a cycle of length at least 4 in $H$, which is incident with at least 4 colors under $\varphi$ since $H$ has already been 4-acyclically colored. Thus, the extended coloring of $G$ is 4-acyclic, a contradiction.

For each cycle $C_{n}$ with $n \geq 4, a_{4}\left(C_{n}\right)=4$, so the bound 4 is sharp.

## 3 4-acyclic colorings of series-parallel graphs

A graph is series-parallel if it has no $K_{4}$-minors. It is known that every series-parallel graph contains a vertex of degree at most 2 Duffin (1965). In this section, we aim to give a sharp upper bound for the 4 -acyclic chromatic number of a series-parallel graph.

Theorem 3.1 For each series-parallel graph $G, a_{4}(G) \leq \Delta(G)+2$ and the bound $\Delta(G)+2$ is sharp.

Proof Suppose that $G$ is a minimal counterexample with the smallest number of $|V(G)|+|E(G)|$. Clearly, $\delta(G) \geq 2$, so by the 2-degeneracy of $G, G$ contains a vertex of degree 2 , say $u$.

For convenience, let $\Delta=\Delta(G)$. Denote the neighbors of $u$ in $G$ by $v$ and $w$. If $v w \in E(G)$, then let $H:=G-\{u\}$; otherwise, let $H:=G-\{u\}+v w$. In any case one can check that $H$ is still a series-parallel graph with $\Delta(H) \leq \Delta$ and thus by the minimality of $G, H$ admits a 4-acyclic $(\Delta+2)$-coloring $\varphi$ with $\varphi(v) \neq \varphi(w)$. Let $S$ be the set of colors on the vertices of $N(v) \cap N(w)$ under $\varphi$. Since $u$ is uncolored under $\varphi$ and $|(N(v) \cap N(w)) \backslash\{u\}| \leq \Delta-1,|S| \leq \Delta-1$. Extend $\varphi$ to a coloring of $G$ by coloring $u$ with a color different from any color in $F:=S \cup\{\varphi(v), \varphi(w)\}$. Since $|F| \leq \Delta+1$, there is one available color for $u$, so the above extension is exercisable. Note that $\varphi(v) \neq \varphi(w)$. By the choice of the color on $u$, one can see that the vertices of every 4 -cycle in $G$ that passes through $u$ are colored all distinctly. Let $C$ be a cycle of length at least 5 in $G$ that passes through $u$. By the definition of $H, C-u+v w$ is a cycle of length at least 4 in $H$, and thus this cycle is incident with at least 4 colors under $\varphi$. This implies that the cycle $C$ in $G$ is also incident with at least 4 colors after the extension of $\varphi$. Therefore, the extended coloring of $G$ is 4 -acyclic, a contradiction.

Since $a_{4}(G)=\Delta(G)+2$ when $G$ is a complete bipartite graph $K_{2, n}$ with $n \geq 2$, the bound $\Delta(G)+2$ in the theorem is sharp.

## $4 r$-acyclic colorings of planar graphs

In this section, we give an upper bound for the $r$-acyclic chromatic number of a planar graph when $r \geq 4$. First, we give some definitions and notations which will be used in our proof. The kthpower $G^{k}$ of a graph $G$ is defined on the same set of vertices as $G$ and has an edge between any pair of vertices of distance at most $k$ in $G$. Agnarsson and Halldórsson (2003) showed the following theorem.

Theorem 4.1 Let $G$ be a planar graph with maximum degree $\Delta$. For any fixed $k \geq$ 1, $G^{k}$ is $O\left(\Delta^{\left\lfloor\frac{k}{2}\right\rfloor}\right)$-colorable. Also, there is a family of graphs that attains this bound.

Clearly, $a_{r}(G) \leq \chi\left(G^{r-1}\right)$. So by Theorem 4.1, we have the following corollary.
Corollary 4.1 Let $G$ be a planar graph with maximum degree $\Delta$. For any fixed $r \geq$ $4, a_{r}(G) \leq C\left(\Delta^{\left\lfloor\frac{r-1}{2}\right\rfloor}\right)$, where $C$ is a constant number.

For sparse planar graphs, we will get an improvement for the bound.

Theorem 4.2 Let $r \geq 4$ and $G$ be a graph with $g(G) \geq 3 r-2$. If every subgraph of $G$ has average degree less than $2+\frac{6}{5 r-5}$, then $a_{r}(G) \leq r$.

Theorem 4.2 gives the following immediate corollary.
Corollary 4.2 If $r \geq 4$ and $G$ is a planar graph with girth at least $\frac{10 r-4}{3}$, or a graph embeddable on the torus or Klein bottle with girth greater than $\frac{10 r-4}{3}$, then $a_{r}(G) \leq r$.

Proof Note that $\frac{10 r-4}{3} \geq 3 r-2$. Clearly every subgraph of $G$ has girth at least as large as the girth of $G$, thus if our conclusion fails, then by Theorem 4.2, the average degree of $G$ is at least $2+\frac{6}{5 r-5}$. It follows that $|E(G)| \geq \frac{5 r-2}{5 r-5}|V(G)|$. We use $v, e$ to denote the number of vertices and edges of $G$, respectively. Let $f$ be the number of faces in an embedding of $G$ on a surface of Euler characteristic $N$ and $g(G)$ be the girth of $G$. By Euler's Formula, we have

$$
2-N=v-e+f \leq e\left(\frac{5 r-5}{5 r-2}-1+\frac{2}{g(G)}\right)=e\left(\frac{2}{g(G)}-\frac{3}{5 r-2}\right) .
$$

Since $N \leq 2$ for the surface mentioned in the corollary, we deduce that $\frac{2}{g(G)} \geq \frac{3}{5 r-2}$, that is, $g(G) \leq \frac{10 r-4}{3}$, and equality holds only when $N=2$. This contradiction completes the proof.

## 5 Proof of Theorem 4.2

Suppose that Theorem 4.2 does not hold. We choose a minimal counterexample $G$ to Theorem 4.2 in terms of $|V(G)|+|E(G)|$. Clearly, $G$ has minimum degree at least two.

A thread in a graph $G$ is a path whose internal vertices have degree 2 in $G$. Two vertices are weak neighbors or weakly adjacent if they are the endpoints of a thread (this includes adjacent vertices, since threads may have no internal vertices). For simplicity, let $[r]=\{1,2, \ldots, r\}$. We give the following claims.

Claim 1 Every thread in $G$ has length at most $r-1$.
Proof of Claim 1 Otherwise, we assume that $G$ has a thread $v_{0} v_{1} \cdots v_{r}$ of length $r$. Consider $H=G-\left\{v_{1}, \cdots, v_{r-1}\right\}$. By our assumption, $a_{r}(H) \leq r$. Suppose that $\pi$ is an $r$-acyclic coloring of $H$ by using the colors in $[r]$. Without loss of generality, assume that $\pi\left(v_{0}\right)=1$. Let $\pi\left(v_{1}\right)=2, \pi\left(v_{2}\right)=3, \ldots, \pi\left(v_{r-1}\right)=r$. If $\pi\left(v_{r-1}\right)=\pi\left(v_{r}\right)$, we recolor $v_{r-1}$ with a color different from $\pi\left(v_{r}\right)$ and $\pi\left(v_{r+2}\right)$, and then we get an $r$-acyclic coloring of $G$, which is a contradiction.

Claim 2 No three vertices of $G$ with degree at least 3 are pairwise weakly adjacent, and no two threads have the same set of endpoints.

Proof of Claim 2 Otherwise, by Claim 1, $G$ has a cycle of length at most $3 r-3$, which contradicts $g(G) \geq 3 r-2$.

When $u$ and $v$ are weakly adjacent, let $l_{u v}$ denote the number of the internal vertices of a shortest $u, v$-thread. (Note that if $u, v$ are adjacent, then $l_{u v}=0$ ). Let $Y=\{v \in$ $V(G): d(v) \geq 3\}$. A weak neighbor $u$ of $v$ is a weak $Y$-neighbor of $v$ if $u \in Y$; Otherwise it is a weak2- neighbor of $v$.

For $v \in V(G)$, let $N_{Y}(v)$ denote the set of weak $Y$-neighbors of $v$ in $G$. For $v \in Y$, let

$$
f(v)=-r+\sum_{u \in N_{Y}(v)}\left(r-l_{v u}-1\right)
$$

By Claim $1, l_{u v} \leq r-2$, so $r-l_{v u}-1 \geq 1$ for each $u \in N_{Y}(G)$. Let $P_{u v}$ denote a $u, v$-thread and $L_{u v}$ denote a sub-thread of $P_{u v}$ by deleting the vertices $u, v$. For any subgraph $A$ of $G$, if $\pi$ is a coloring of $A$, then let $C(A)=\{\pi(v) \mid v \in A\}$. A vertex coloring of $A$ such that any two vertices of $A$ have distinct colors is called a rainbow coloring of $A$.

Claim 3 If $v \in Y$, then $f(v) \geq 1$.
Proof of Claim 3 Clearly, if for some $u \in N_{Y}(v), l_{u v}=0$, then $f(v) \geq 1$, so we assume that $l_{u v} \geq 1$ for $u \in N_{Y}(v)$. Let $H$ be the graph obtained from $G$ by deleting $v$ and all its weak 2-neighbors. By our assumption, $a_{r}(H) \leq r$. Let $\pi$ be an $r$ acyclic coloring of $H$ by using the colors in [r]. Assume $d(v)=m$ and $N_{Y}(v)=$ $\left\{u_{1}, u_{2}, \cdots, u_{m}\right\}$.

Now let $\pi(v) \in[r] \backslash \pi\left(u_{1}\right)$. If $\sum_{u \in N_{Y}(v)}\left(r-l_{u v}-1\right)-1 \leq r-1$, then there exist $S_{1}, S_{2}, \cdots, S_{m}$ such that $\pi\left(u_{1}\right) \notin S_{1}$ and $S_{1} \cup S_{2} \cup \cdots S_{m} \subseteq[r] \backslash \pi(v)$, where $\left|S_{1}\right|=r-l_{u_{1} v}-2,\left|S_{i}\right|=r-l_{u_{i} v}-1$ for $2 \leq i \leq m$ and $S_{i} \cap S_{j}=\phi$, for $i \neq j$. Now we give each thread $L_{u_{i v}}$ a rainbow coloring using the colors in $[r] \backslash\left(\pi(v) \cup S_{i}\right)$ for $i=2, \ldots, m$. We then give $L_{u_{1 v}}$ a rainbow matching using the colors in $[r] \backslash(\pi(v) \cup$ $\left.\pi\left(u_{1}\right) \cup S_{1}\right)$. For each thread $P\left(u_{i v}\right)$, if $\pi\left(u_{i}\right)=\pi(x)$, where $x$ is the neighbor of $u_{i}$ in thread $P_{u_{i v}}$, then we recolor $x$ to obtain a proper coloring of $P_{u_{i}}$. It follows that $C\left(P_{u_{i v}}\right) \cup C\left(P_{u_{j v}}\right)=[r]$, for $i \neq j$. Thus we get an $r$-acyclic coloring of $G$, which is a contradiction, so $\sum_{u \in N_{Y}(v)}\left(r-l_{u v}-1\right)-1 \geq r$ and thus $f(v) \geq 1$.
Claim 4 If $v \in Y$, then $\sum_{u \in N_{Y}(v)} f(u) \geq r+2$.
Proof of Claim 4 Suppose on the contrary, $\sum_{u \in N_{Y}(v)} f(u) \leq r+1$. For a vertex $u \in N_{Y}(v)$, if it satisfies that $f(u) \leq r-l_{u v}-1$, then we call it $v$-good, otherwise we call it $v-b a d$. For convenience, let $N_{Y}^{g}(v)$ denote the set of $v-\operatorname{good}$ vertices and $N_{Y}^{b}(v)$ denote the set of $v-b a d$ vertices. Note that when $u$ is $v-b a d$, then $l_{u v} \geq 1$, since otherwise we have that $f(u) \geq r-l_{u v}=r$, hence $\sum_{u \in N_{Y}(v)} f(u) \geq r+2$, which is a contradiction. Let $H$ be the graph obtained from $G$ by deleting the vertex $v$, the $v$-good neighbors, and all their weak 2-neighbors. Let $H_{1}$ denote the subgraph induced by $V-V(H)$. By assumption, $H$ has an $r$-acyclic coloring $\pi$. First we color $v$ such that $\pi(v)=r$.

Suppose that $u$ is $v$-good. Consider each thread from $u$ except $u, v$-thread. Let

$$
s(u)=r-\sum_{w \in N_{Y}(u), w \neq v}\left(r-l_{u w}-1\right)+1 .
$$

We then have

$$
\begin{aligned}
s(u) & =1+r-\sum_{w \in N_{Y}(u)}\left(r-1-l_{u w}\right)+\left(r-1-l_{u v}\right) \\
& =1-f(u)+r-1-l_{u v} \\
& \leq 1 .
\end{aligned}
$$

Fact For any subset $S(u) \subseteq[r-1]$ of size $s(u)$, we have a proper coloring of each thread $P_{u w}\left(w \in N_{Y}(v)-v\right)$ such that
A. $S(u) \subseteq C\left(P_{u w}\right)$;
B. $C\left(P_{u w_{i}} \cup P_{u w_{j}}\right)=[r]$, for $w_{i}, w_{j} \in N_{Y}(u)-v$ and $i \neq j$.

Proof of Fact First suppose that there is a thread (say $L_{u w_{1}}$, where $w_{1} \in N_{Y}(u)-v$ ) such that $l_{u w_{1}}=0$, then we have that $s(u)=1, d(u)=3$ and $l_{u w_{2}}=r-2$. Let $S(u)=\{c\}$. If $\pi\left(w_{1}\right) \neq c$, then let $\pi(u)=c$ and give $L_{u w_{2}}$ a rainbow coloring by using the colors in $[r] \backslash\left(\pi\left(w_{1}\right) \cup \pi(u)\right)$. In thread $L_{u w_{2}}$, if the neighbor of $w_{2}$ has the same color with $w_{2}$, we change it to get a proper coloring of $P_{u w_{2}}$. If $\pi\left(w_{1}\right)=c$, we choose a color from $[r-1] \backslash\left(\pi\left(w_{1}\right) \cup \pi\left(w_{2}\right)\right)$ for $u$, then give a rainbow coloring to $L_{u w_{2}}$ by using the colors from $[r] \backslash\left(\pi(u) \cup \pi\left(w_{2}\right)\right)$. Such coloring satisfies $A$ and $B$, so in the following we assume that $l_{u w} \geq 1$ for $w \in N_{Y}(u)-v$.

If $s(u)=1$ and $C\left(N_{Y}(u)-v\right)=S(u)$, then we color $u$ by a color in $[r-1] \backslash S(u)$ and give each thread $L_{w u}$ for $w \in N_{Y}(u)-v$ a rainbow coloring using the colors in $[r] \backslash(S(u) \cup \pi(u))$. Let $T=[r] \backslash(S(u) \cup \pi(u))$ and let $M\left(L_{u w}\right)=T \backslash C\left(L_{u w}\right)$ denote the colors which are in $T$ and do not appear in thread $L_{u w}$. Clearly $\left|M\left(L_{u w}\right)\right|=$ $r-2-l_{u w}$. Since $\sum_{w \in N_{Y}(u)-v}\left(r-l_{u w}-2\right) \leq r-2=r-s(u)-1=|T|$, we further assume that $M\left(L_{u w_{i}}\right) \cap M\left(L_{u w_{j}}\right)=\emptyset$, for $w_{i}, w_{j} \in N_{Y}(u)-v$ with $i \neq j$. It follows that $C\left(P_{u w_{i}} \cup P_{u w_{j}}\right)=[r]$, for $w_{i}, w_{j} \in N_{Y}(u)-v$ and $i \neq j$. Finally, in each thread $P_{u w}$ for $w \in N_{Y}(u)-v$, we recolor the neighbor of $w$ if necessary to get a proper coloring of $P_{u w}$. Now it is easy to see this coloring satisfies $A$ and $B$.

Now we assume that $s(u) \geq 2$ or $C\left(N_{Y}(u)-v\right) \neq S(u)$. We color $u$ such that $\pi(u) \in S(u)$ and $\pi(u) \neq \pi(w)$ for some $w \in N_{Y}(u)-v$. Without loss of generality, we assume that $\pi(u) \neq \pi\left(w_{1}\right)$. Let $T=[r] \backslash S(u)$. We give a rainbow coloring to each thread $L_{u w}\left(w \notin\left\{w_{1}, v\right\}\right)$ using the colors in $[r] \backslash \pi(u)$ such that $(S(u) \backslash \pi(u)) \subseteq C\left(L_{u w}\right)$. If $\pi\left(w_{1}\right) \in S(u)$, we give a rainbow coloring to $L_{u w_{1}}$ using the colors in $[r] \backslash\left\{\pi(u), \pi\left(w_{1}\right)\right\}$ such that the colors in $S(u) \backslash\left(\pi\left(w_{1}\right) \cup S(u)\right)$ appear on $L_{u w_{1}}$. If $\pi\left(w_{1}\right) \notin S(u)$, we give a rainbow coloring to $L_{u w_{1}}$ using the colors in $[r] \backslash\left\{\pi(u), \pi\left(w_{1}\right)\right\}$ such that $(S(u) \backslash \pi(u)) \subseteq C\left(L_{u w_{1}}\right)$. Next we give a rainbow coloring to each thread $L_{u w}\left(w \in N_{Y}(u)-\left\{w_{1}, v\right\}\right)$ by using the colors in $[r] \backslash \pi(u)$ such that $(S(u) \backslash \pi(u)) \subseteq C\left(L_{u w}\right)$.

Let $M\left(L_{u w}\right)=T \backslash C\left(L_{u w}\right)$. If $\pi\left(w_{1}\right) \in S(u)$, then $\left|M\left(L_{u w_{1}}\right)\right|=r-s(u)-\left(l_{u w_{1}}-\right.$ $s(u)+2)=r-l_{u w_{1}}-2$. If $\pi\left(w_{1}\right) \notin S(u)$, then $\left|M\left(L_{u w_{1}}\right)\right|=r-s(u)-\left(l_{u w_{1}}-\right.$ $s(u)+1)=r-l_{u w_{1}}-1$. For $w \in N_{Y}(u)-\left\{w_{1}, v\right\}$, it holds that $\left|M\left(L_{u w}\right)\right|=$ $r-s(u)-\left(l_{u w}-s(u)+1\right)=r-l_{u w}-1$. Recall that $\sum_{w \in N_{Y}(u)-\{v\}}\left(r-l_{u w}-1\right)=$ $r-s(u)+1=|T|+1$. If $\pi\left(w_{1}\right) \in S(u)$, we may assume that $M\left(L_{u w_{i}}\right) \cap M\left(P_{u w_{j}}\right)=\emptyset$ for $i \neq j$. If $\pi\left(w_{1}\right) \notin S(u)$, we may assume that $M\left(L_{u w_{i}}\right) \cap M\left(P_{u w_{j}}\right)=\emptyset$ for $i \neq j$
except $M_{R}\left(L_{u w_{1}}\right) \cap M\left(P_{u w_{2}}\right)=\pi\left(w_{1}\right)$, so $C\left(P_{u w_{i}} \cup P_{u w_{j}}\right)=[r]$, for $i \neq j$. At last, in each thread $P_{u w}\left(w \in N_{Y}(u)-v\right)$, we change the neighbor of $w$ if necessary to get a proper coloring of $P_{u w}$. Hence $A$ and $B$ hold. This completes the proof of the fact.

$$
\begin{aligned}
& \text { Assume that } d(v)=m \geq 3 \text { and } N_{Y}(v)=\left\{u_{1}, u_{2}, \cdots, u_{m}\right\} \text {. Since } \sum_{u \in N_{Y}(v)} f(u) \leq r, \\
& \qquad \begin{aligned}
& \sum_{u \in N_{Y}^{g}(v)}\left(r-1-s(u)-l_{u v}\right)+\sum_{u \in N_{Y}^{b}(v)}\left(r-1-l_{u v}\right) \\
& \leq \sum_{u \in N_{Y}^{g}(v)}(f(u)-1)+\sum_{u \in N_{Y}^{b}(v)}(f(u)-1) \\
= & \sum_{u \in N_{Y}(v)} f(u)-d(u) \\
\leq & r-1 .
\end{aligned}
\end{aligned}
$$

Thus, there exist $S_{1}, S_{2}, \cdots, S_{m}$ such that $S_{1} \cup S_{2} \cup \cdots S_{m} \subseteq[r-1]$, where $\left|S_{i}\right|=$ $r-s\left(u_{i}\right)-l_{u_{i v}}-1$ if $u_{i}$ is $v$-good and $\left|S_{i}\right|=r-l_{u_{i v}}-1$ if $u_{i}$ is $v$-bad. Moreover, we assume that $S_{i} \cap S_{j}=\phi$, for $i \neq j$.

If $u_{i} \in N_{Y}(v)$ is $v$-bad, then we give a rainbow coloring of $L_{u_{i v}}$ using the colors in $[r-1] \backslash S_{i}$. If $u_{i}$ is $v$-good, then we choose $S\left(u_{i}\right) \in[r-1] \backslash S_{i}$ such that $\left|S\left(u_{i}\right)\right|=s\left(u_{i}\right)$. By the above fact, we can properly color each thread of $P_{u_{i w}}\left(w \in N_{Y}\left(u_{i}\right)-v\right)$ such that $S\left(u_{i}\right) \subseteq C\left(P_{u i w}\right)$ and $C\left(P_{u_{i w j}} \cup P_{u_{i w k}}\right)=[r]$, for $w_{k}, w_{j} \in N_{Y}\left(u_{i}\right)-v$ and $w_{k} \neq w_{j}$. Give a rainbow coloring of $L_{u_{i v}}$ using the colors in $[r-1] \backslash S_{i} \cup S\left(u_{i}\right)$. In each thread $P_{u_{i v}}$, we adjust the color of the neighbor of $u_{i}$ if necessary to get a proper coloring of $P_{u_{i v}}$. Thus, we have a proper coloring of $H_{1}$, moreover for any two vertices $x, y \in V\left(H_{1}\right)$ and any path $Q(x, y)$ between $x$ and $y$ in $H_{1}, C(Q(x, y))=$ [ $r$ ]. Thus, we get an $r$-acyclic coloring of $G$, which is a contradiction. Therefore, $\sum_{u \in N_{Y}(v)} f(u) \geq r+2$. This completes the proof of Claim 4.

We complete the proof using discharging method. Let $d(v)$ be the initial charge on the vertex $v \in V(G)$. We move charge from vertex to vertex, without changing the total according to the following rules:
a. Every $v \in Y$ gives each weak 2-neighbor the amount $\frac{3}{5 r-1}$.
b. Every $v \in Y$ gives each weak $Y$-neighbor the amount $\frac{3 f(v)+(r+2)(d(v)-3)}{(5 r-1) d(v)}$.

Claim 5 Every $v \in Y$ receives from its weak $Y$-neighbors at least $\frac{r+2}{5 r-1}$.
Proof of Claim 5 If every $u \in N_{Y}(v)$ sends $v$ at least $\frac{f(u)}{5 r-1}$, then $v$ receives from $N_{Y}(v)$ at least $\frac{1}{5 r-1} \sum_{u \in N_{Y}(v)} f(u) \geq \frac{r+2}{5 r-1}$, by Claim 4.

Otherwise, for some $u \in N_{Y}(v)$, it holds that

$$
\frac{3 f(u)+(r+2)(d(u)-3)}{(5 r-1) d(u)}<\frac{f(u)}{5 r-1},
$$

that is,

$$
(r+2)(d(u)-3)<f(u)(d(u)-3),
$$

so we conclude that $d(u) \geq 4$ and $f(u)>r+2$. Thus, $u$ gives to $v$ at least

$$
\frac{3 f(u)+(r+2)(d(u)-3)}{(5 r-1) d(u)} \geq \frac{3(r+2)+(r+2)(d(u)-3)}{(5 r-1) d(u)}=\frac{r+2}{5 r-1} .
$$

Moreover, all other amounts to $v$ are nonnegative, since if $y \in N_{Y}(v)$, then $d(w) \geq$ 3 and $f(w) \geq 1$.

Let $\hat{d}(v)$ denote the new charge of $v$ after discharging.
Claim 6 After the discharging, it holds that $\hat{d}(v) \geq 2+\frac{4 d(v)-2}{5 r-1}$, for all $v \in V(G)$.
Proof of Claim 6 If $d(v)=2$, then $v$ sends out zero and receives $\frac{3}{5 r-1}$ from each of its two weak $Y$-neighbors, so $\hat{d}(v)=2+\frac{6}{5 r-1}=2+\frac{4 d(v)-2}{5 r-1}$.

Now consider $v \in Y$. By the discharging rule, $v$ sends out $\frac{3}{5 r-1} \sum_{w \in N_{Y}(v)} l_{v w}$ to its weak 2-neighbors and $\frac{3 f(v)+(r+2)(d(v)-3)}{5 r-1}$ to its weak $Y$-neighbors. By Claim 5, $v$ receives at least $\frac{r+2}{5 r-1}$ from its weak $Y$-neighbors, so

$$
\begin{aligned}
\hat{d}(v) & \geq d(v)-\frac{3}{5 r-1} \sum_{w \in N_{Y}(v)} l_{w v}-\frac{3 f(v)+(r+2)(d(v)-3)}{5 r-1}+\frac{r+2}{5 r-1} \\
& =d(v)-\frac{3}{5 r-1}\left[-r+\sum_{w \in N_{Y}(v)}\left(r-l_{w v}+l_{w v}-1\right)\right]-\frac{(r+2)(d(v)-4)}{5 r-1} \\
& =\frac{r d(v)+7 r+8}{5 r-1}
\end{aligned}
$$

Since $d(v) \geq 3$ and $r \geq 4$, we have

$$
r d(v)+7 r+8=(d(v)-3) r+10+10 r-2 \geq 4 d(v)-2+10 r-2
$$

Therefore,

$$
\frac{(r+2) d(v)+7 r+4}{5 r-1} \geq 2+\frac{4 d(v)-2}{5 r-1}
$$

and the proof of Claim 6 completes.
Now we have that $\hat{d}(v) \geq 2+\frac{4 d(v)-2}{5 r-1}$, for all $v \in V(G)$. It follows that

$$
\begin{aligned}
2|E(G)| & =\sum_{v \in V(G)} \hat{d}(v) \geq \sum_{v \in V(G)}\left(2+\frac{4 d(v)-2}{5 r-1}\right) \\
& =2\left(1-\frac{1}{5 r-1}\right)|V(G)|+\frac{8}{5 r-1}|E(G)|,
\end{aligned}
$$

and hence

$$
\frac{5 r-2}{5 r-1}|V(G)| \leq \frac{5 r-5}{5 r-1}|E(G)| .
$$

Thus, the average degree of $G$ is at least $2+\frac{6}{5 r-5}$, which gives a contradiction. This completes the proof of Theorem 4.2.

## 6 Remark

Let $r \geq 4$ be an integer. We propose the following problems for further research.
Problem 1 What is the best upper bound for $a_{r}(G)$ when $G$ is a planar graph ?
Problem 2 What is the best upper bound for $a_{r}(G)$ when $G$ is a planar graph containing no copy of $K_{2, n}$ or even no $C_{4}$ ?

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