

Principles of Communications

Chapter IV: Digitization of Analog Signal

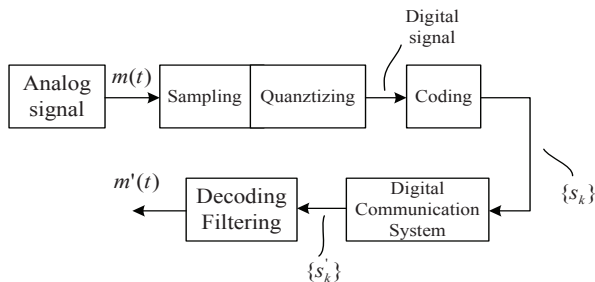
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4.1 Introduction



In the nature, most signals, such as audio, image, are analog. One wants to transmit analog signal using digital techniques.

Analog to Digital

- ▶ Step 1. Sampling: from the continuous to the discrete.
- ▶ Step 2. Quantization: digital signal.
- ▶ Step 3. Code: data compression.

4.2 Sampling of Analog Signal Low-Pass Signal

4.2.1 Sampling of Low-Pass Signal

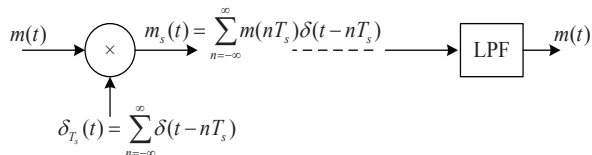
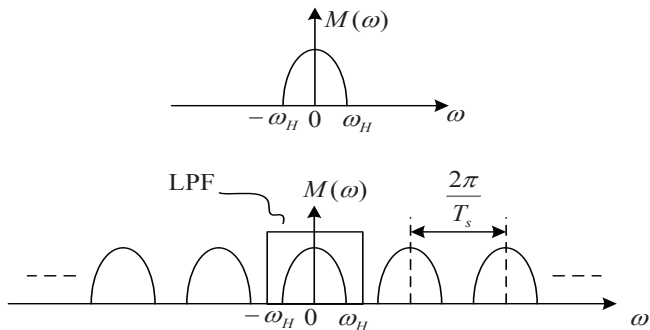


Figure: Scheme of sampling of analog signal

Low-pass sampling theorem:

If the highest frequency of a continuous analog signal $m(t)$ is less than f_H , and if its sampled by interval time $T_s \leq 1/2f_H$, then $m(t)$ can be completely decided by these samples.



Discussions

- ▶ $T_s \leq 1/2f_H \Leftrightarrow f_s \geq 2f_H$.
- ▶ $m(t) = \sum_{n=-\infty}^{\infty} m(nT_s) \text{sinc} \omega_H(t - nT_s)$. One can recover $m(t)$ utilizing $m(nT_s)$ through filtering.
- ▶ How about $f_s \leq 2f_H$? Aliasing? Can you draw a picture?

4.2.2 Sampling of Band-Pass Signal

A sufficient no-loss condition for sampling signals that do not have baseband components exists that involves the width of the non-zero frequency interval as opposed to its highest frequency component. The sufficient condition is that the frequency band of $m(t)$ is within the range

$$\left(\frac{N}{2} f_s, \frac{N+1}{2} f_s \right),$$

where N is some nonnegative integer. The reconstruction filter is an ideal bandpass filter with cutoffs at the upper and lower edges of the specified band, which is the difference between a pair of low-pass impulse responses

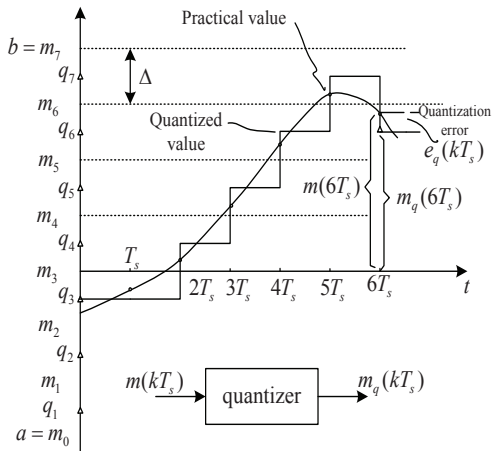
$$(N+1)f_s \operatorname{sinc}((N+1)\pi f_s t) - Nf_s \operatorname{sinc}(N\pi f_s t).$$

See details in the webpage:

<http://en.wikipedia.org/wiki/Sampling-theorem>.

4.3 Quantization of Sampled Signal

1 Uniform Quantization



Uniform quantization (quantizer parameters)

- ▶ the input max/min values are a and b , respectively.
- ▶ the number of quantization levels is M .
- ▶ the i th quantization interval is (m_{i-1}, m_i) . Then quantization interval $\Delta_i = \Delta = (b - a)/M$ and the quantized output level $q_i = (m_{i-1} + m_i)/2$.

- ▶ Quantization error e_q :
 - ▶ If the sample $m(kT_s) \in [a, b]$, quantization error $e_q \leq \Delta/2$.
 - ▶ If the sample $m(kT_s) \notin [a, b]$ (overload area), quantization error e_q may be larger than $\Delta/2$.
 - ▶ When designing the quantizer, we always avoid the overload phenomenon.
- ▶ Quantization noise: Assume that $m(t)$ is stable random process with zero expect value and probability density function $f(x)$, and $m(t) \in [a, b]$. The quantization noise

$$N_q = E[|m(kT_s) - m_q(kT_s)|^2] = \sum_{i=1}^M \int_{m_{i-1}}^{m_i} (x - q_i)^2 f(x) dx.$$

- ▶ Quantization signal-to-noise ratio:

$$\frac{S}{N_q} = \frac{E[|m(kT_s)|^2]}{E[|m(kT_s) - m_q(kT_s)|^2]}$$

- ▶ Example. The number of quantization levels is M , the pdf of $m(t) U[-a, a]$. Determine the $\frac{S}{N_q}$?

▶

$$\begin{aligned} N_q &= \sum_{i=1}^M \int_{m_{i-1}}^{m_i} (x - q_i)^2 \frac{1}{2a} dx \\ &= \sum_{i=1}^M \int_{-a+(i-1)\Delta}^{-a+i\Delta} \left(x + a - i\Delta + \frac{\Delta}{2}\right)^2 \frac{1}{2a} dx \\ &= \sum_{i=1}^M \frac{1}{2a} \frac{\Delta^3}{12} \\ &= \frac{(\Delta)^2}{12} \quad (\text{since } M\Delta = 2a). \end{aligned}$$



$$\begin{aligned} S &= E[m(kT_s)^2] = \int_{-a}^a x^2 \frac{1}{2a} dx \\ &= \frac{a^2}{3} = \frac{M^2}{12} \Delta^2 \end{aligned}$$

- ▶ Result: $(\frac{S}{N_q})_{dB} = 20 \log M = 6N$. Usually, we say, if the quantization bits are increased by 1, the $(\frac{S}{N_q})_{dB}$ increases 6dB.
- ▶ Drawbacks: uniform quantizer is not good for small signals. However, in the nature, such as speech signal, most cases are small.

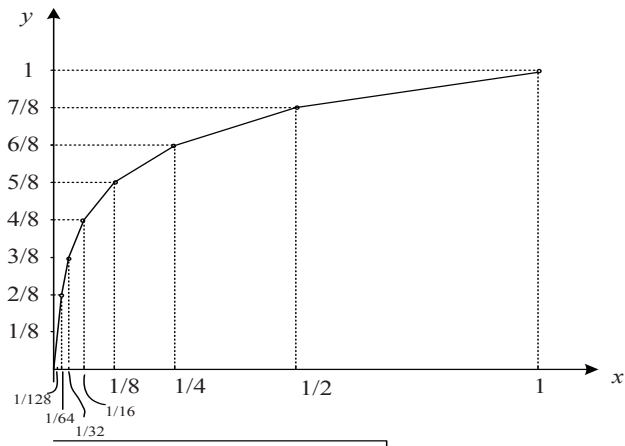
4.3.3 Non-uniform Quantization

In the non-uniform quantization, the interval of quantization varies with samples of a signal. It means that the response, $y = f(x)$, of the quantizer is nonlinear. In practice, ITU recommend two kinds of quantizer:

- ▶ A-law compression: $y = \begin{cases} \frac{Ax}{1+\ln A}, & 0 < x \leq \frac{1}{A}, \\ \frac{1+\ln Ax}{1+\ln A}, & \frac{1}{A} \leq x \leq 1. \end{cases}$
- ▶ μ -law compression: $y = \frac{\ln(1+\mu x)}{\ln(1+\mu)}, 0 \leq x \leq 1.$

In America and Japan, μ -law 15 broke line compression $\mu = 255$ is exploited; In our country and Europe, A-law 13 broke line compression ($A=87.6$) is exploited.

A-law Broke Line Compression Characteristics



- ▶ The above characteristics is named A-law compression when $A=87.6$. See the pp.95, table 4.3.2.
- ▶ Why is 13 broke line?

4.4 Pulse Code Modulation - PCM

Table: Folded binary code and nature binary code

Polarity	Nature	Folded	Levels
+	1111	1111	15
	1110	1110	14
	1101	1101	13
	1100	1100	12
	1011	1011	11
	1010	1010	10
	1001	1001	9
	1000	1000	8
-	0111	0000	7
	0110	0001	6
	0101	0010	5
	0100	0011	4
	0011	0100	3
	0010	0101	2
	0001	0110	1
	0000	0111	0

- Folded binary code is good for small signals.

4.4.1 Principles of Pulse Code Modulation (PCM)

The PCM consists of three steps: sampling, quantization, and **coding**.

- ▶ A-law 13 broke line: the input range $[0, 1]$ is non-uniformly divided into 8 segments, $[0, 1/128]$, ..., $[1/4, 1/2]$, $[1/2, 1]$.
- ▶ Each segment is uniformly divided into 16 smaller segments.

Coding rules: Folded binary code

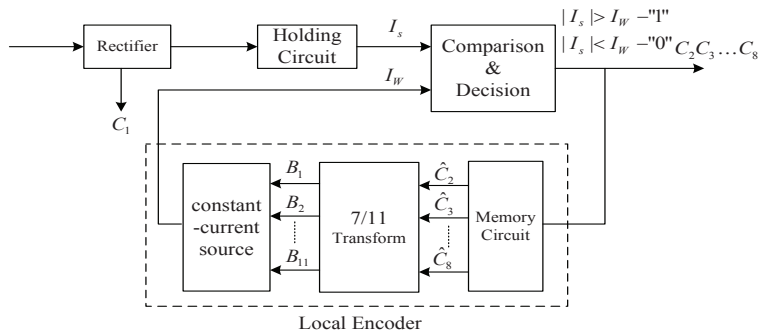
- ▶ Polarity bit, named c_1 , to denote the positive or negative of the pulse, e.g., positive – $c_1 = 1$ and negative – $c_1 = 0$.
- ▶ Segment bits, named $c_2c_3c_4$, to denote which segment the pulse locates.
- ▶ Inner-segment bits, named $c_5c_6c_7c_8$, to denote which quantization level the pulse is.

Define the minimum quantization interval $\Delta = \frac{1}{2048}$ as the quantization unit. Then, we have the table

No. of segment	beginning level	quantization interval
1 000	0	Δ
2 001	16Δ	Δ
3 010	32Δ	2Δ
4 011	64Δ	4Δ
5 100	128Δ	8Δ
6 101	256Δ	16Δ
7 110	512Δ	32Δ
8 111	1024Δ	64Δ

- ▶ Exe: pulse $I_s = 1260\Delta$, determine the A-law 13 broke line code.

PCM Encoder

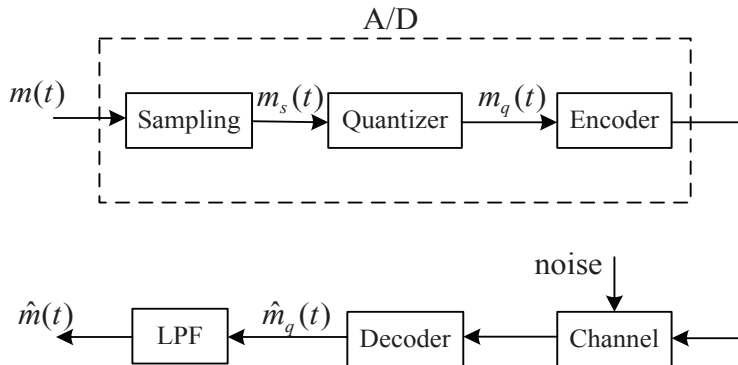


To determine c_i , set $\hat{c}_i=1$

Encoder Rate

- ▶ Assume the speech signal $[0, 4kHz]$.
- ▶ Sampling rate $f_s = 8k$ samples/s.
- ▶ PCM A-law 13 broke line rule: $R_b = 64 kbit/s$.

Quantization Noise in PCM System



- ▶ The output of the LPF: $\hat{m}(t) = m(t) + n_q(t) + n_e(t)$.
- ▶ The output SNR: $\frac{S_o}{N_o} = \frac{E[m^2(t)]}{E[n_q^2(t) + E[n_e^2(t)]}$.
- ▶ Assume the input signal $m(t) \sim U[-a, a]$ and quantizer is with N bit (M levels) uniform quantization. Determine $\frac{S_o}{N_o}$?

- ▶ Signal to quantization noise ratio:

$$\frac{S_o}{N_q} = \frac{E[m^2(t)]}{n_q^2(t)} = M^2 = 2^{2N}.$$

- ▶ Signal to noise ratio (error bits):

- ▶ Noise power (error bits):

$$N_e = E[n_e^2(t)] = P_e \sum_{i=1}^N (2^{i-1} \Delta)^2 = \frac{2^{2N} - 1}{3} \Delta^2 P_e \approx \frac{2^{2N}}{3} \Delta^2 P_e$$

- ▶ Signal power: $S_o = E[m^2(t)] = \int_{-a}^a x^2 \frac{1}{2a} dx = \frac{a^2}{3} = \frac{2^{2N}}{12} \Delta^2.$

$$(2a = 2^N \Delta)$$

- ▶ We have $\frac{S_o}{N_e} = \frac{1}{4P_e}.$

- ▶ We obtain the output SNR: $\frac{S_o}{N_o} = \frac{1}{\frac{N_q}{S_o} + \frac{N_e}{S_o}} = \frac{2^{2N}}{1 + 4P_e 2^{2N}}.$

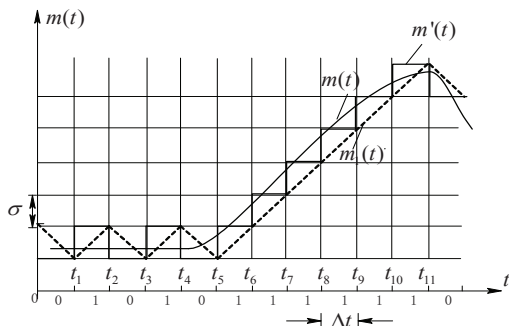
- ▶ Discussions: (From the viewpoint of the practice)

- ▶ High SNR: P_e is very small, $4P_e 2^{2N} \ll 1$, $\frac{S_o}{N_o} \approx \frac{S_o}{N_q}.$
- ▶ Low SNR: P_e is large, $4P_e 2^{2N} \gg 1$, $\frac{S_o}{N_o} \approx \frac{S_o}{N_e}.$

4.6 Delta Modulation

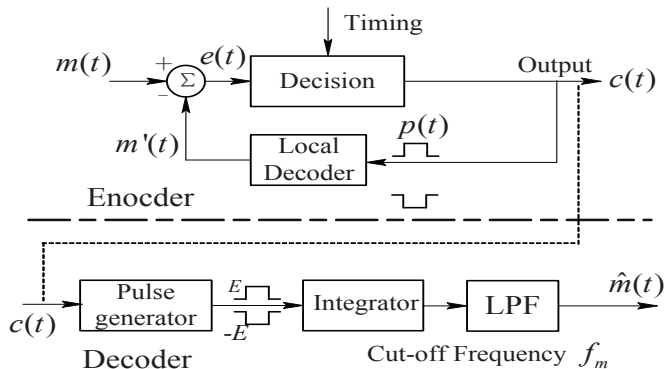
1. Principles of Delta Modulation

Definition: ΔM is to encode the difference between the neighboring samples.



- ▶ In the above figure, if Δt and σ are small enough, staircase waveform, $m'(t)$, can be arbitrarily close to $m(t)$.
- ▶ Encoding: “1” — $m'(t)$ goes upstairs δ and “0” — $m'(t)$ goes downstairs δ , $m'(t)$ is equivalent to binary sequence.

2. Encoder and Decoder



- Decision Rules:

$$m(t) - m'(t)|_{t=t_-} \begin{cases} > 0, & \text{"1"} \\ < 0, & \text{"0"}. \end{cases}$$

3. Overload Characteristic

- ▶ Quantization noise: $e_q(t) = m(t) - m'(t)$
 - ▶ Normal quantization noise: $|e_q(t)| \leq \sigma$.
 - ▶ Overload quantization noise: If the output $m'(t)$ of the local encoder cannot pace up with variation of $m(t)$, $|e_q(t)| > \sigma/2$.
- ▶ Maximum tracking slope:
 - ▶ Assume sampling interval Δt — sampling rate $f_s = \frac{1}{\Delta t}$;
 - ▶ Quantization step: σ .
 - ▶ Maximum tracking slope: $K = \sigma f_s$.
- ▶ To avoid overload quantization noise:

$$\left| \frac{dm(t)}{dt} \right| \leq \sigma f_s$$

- ▶ To elevate the tracking ability: increase quantization step σ and sampling rate f_s .
- ▶ Drawbacks: Increase quantization noise and transmission rate.

4. Coding Range and Signal to Quantization SNR

Maximum coding level and Minimum coding level

- ▶ Maximum coding level: If input signal $m(t) = A \sin \omega_k t$, the maximum slope $K = A\omega_k$. To avoid overload noise, $A\omega_k \leq \sigma f_s$ is desired. So $A_{\max} = \frac{\sigma f_s}{2\pi f_k}$.
- ▶ Minimum coding level: If $m(t) \in \left[-\frac{\sigma}{2}, \frac{\sigma}{2}\right]$, the output coded sequence is 10101010 \dots and cannot show the characteristic of the signal. So $A_{\min} = \frac{\sigma}{2}$.

Assume the input signal $m(t) = A \sin \omega_k t$. Quantization SNR

$$\begin{aligned}\frac{S_o}{N_q} &= \frac{3}{8\pi^2} \left(\frac{f_s^3}{f_k^2 f_m} \right) \approx 0.04 \frac{f_s^3}{f_k^2 f_m} \\ &= 30 \log f_s - 20 \log f_k - 10 \log f_m - 14(\text{dB}).\end{aligned}$$

where f_m is the cut-off frequency of the low pass filter.

- ▶ $f_s \rightarrow 2f_s$, $\left(\frac{S_o}{N_q}\right)_{\text{dB}} \uparrow 9\text{dB}$; $f_k \rightarrow 2f_k$, $\left(\frac{S_o}{N_q}\right)_{\text{dB}} \downarrow 6\text{dB}$;
 $f_m \rightarrow 2f_m$, $\left(\frac{S_o}{N_q}\right)_{\text{dB}} \downarrow 3\text{dB}$;

5. Comparisons of PCM and ΔM

	PCM	ΔM
Sampling Rate	$f \geq 2f_m$	related to the maximum tracking slope, usually larger than Nyquist rate.
Bit rate	$\geq 64kbit/s$	determined by f_s
$\frac{S_o}{N_q}$	$6N_{dB}$	$\frac{S_o}{N_q} = \frac{3}{8\pi^2} \left(\frac{f_s^2}{f_k^2 f_m} \right)$
Effects of error bits	sensitive	insensitive, applied to the system with high BER
Complexity	high, but widely used	low and only used in special applications. (why?)

Assume PCM and ΔM have the same transmission rate, R_b .

▶ ΔM : $f_s = R_b$.

▶ PCM: $f_s = 2f_m = \frac{R_b}{N} \rightarrow R_b = 2Nf_m$

$$\frac{S_o}{N_q} \approx 10 \log \left(0.04 \frac{f_s^3}{f_k^2 f_m} \right)$$

$$= 10 \log \left(0.32 N^3 \frac{f_m^2}{f_k^2} \right)$$

When $f_k = 1kHz$ and

$f_m = 4kHz$,

$$\frac{S_o}{N_q} \approx 30 \log 1.72N \text{ dB.}$$

